

Analysis of FOX

Jorge Nakahara Jr

jorge_nakahara@yahoo.com.br

Presented by: Mads Rasmussen

mads@sitec.org.br

SBSeg 2008, Gramado, RS, Brazil

Summary

- The FOX block cipher family
- Information leakage in FOX
- Impossible-Differential Analysis
- Conclusions

SBSeg 2008, Gramado, RS, Brazil

The FOX Cipher Family

- **Designers: P. Junod and S. Vaudenay (2004)**
- **New name: IDEA-NXT**
- **Owned by Mediacrypt A.G.**
- **Lai-Massey structure, like IDEA cipher**
- **Two main ciphers: FOX64 and FOX128**
- **Byte-oriented design: byte are elements of $GF(2^8) = GF(2)[x]/(x^8+x^7+x^6+x^5+x^4+x^3+1)$**

SBSeg 2008, Gramado, RS, Brazil

The FOX Cipher Family

- Each round of FOX64 uses a bijective mapping called f32 consisting of an MDS matrix, xor with two round subkeys, and an 8x8-bit S-box
- Each round of FOX128 uses a bijective mapping called f64, consisting of an MDS matrix, xor with two round subkeys and an 8x8-bit S-box
- To break the round symmetry, there is a linear transformation called orthomorphism (**or**) at the end of each round.

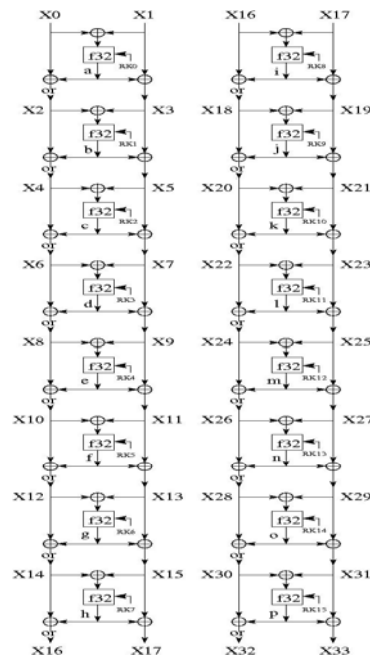
SBSeg 2008, Gramado, RS, Brazil

The FOX Cipher Family

| Cipher | Block Size (bits) | Key Size (bits) | #rounds |
|-------------------|-------------------|------------------------------|-------------------------------|
| FOX64 | 64 | 128 | 16 |
| FOX128 | 128 | 256 | 16 |
| FOX64/k/r | 64 | k ($0 \leq k \leq 256$) | r ($12 \leq r \leq 255$) |
| FOX128/k/r | 128 | k ($0 \leq k \leq 256$) | r ($12 \leq r \leq 255$) |

SBSeg 2008, Gramado, RS, Brazil

FOX64



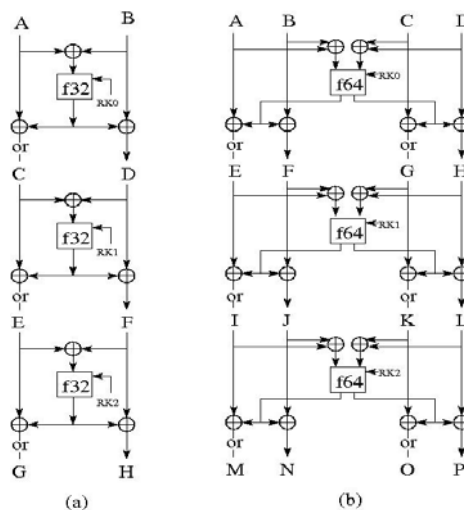
SBSeg 2008, Gramado, RS, Brazil

Information Leakage

- **Information leakage: discover (non-trivial) internal cipher data using only (known) plaintext-ciphertext pairs.**
- **More specifically, for FOX, use the high level Lai-Massey scheme.**
- **Such kind of leakage absent in IDEA cipher.**
- **Leakage leads to full break of FOX64/k/2 and FOX128/k/2.**

SBSeg 2008, Gramado, RS, Brazil

3-round FOX64 and FOX128



SBSeg 2008, Gramado, RS, Brazil

Information Leakage

- Consider 2-round FOX64, for ex., the top 2 rounds of the previous slide.
- From the Lai-Massey scheme:
 $A \oplus B = \text{io}(C) \oplus D$, $C \oplus D = \text{io}(E) \oplus F$, where \oplus is exclusivo-or and **io** is the inverse of **or**.
- Known-plaintext (KP) setting.
- Thus, $D \oplus \text{or}(D) = \text{or}(A \oplus B) \oplus \text{io}(E) \oplus F$
 and $C \oplus \text{io}(C) = A \oplus B \oplus \text{io}(E) \oplus F$, and
 both C and D can be uniquely determined.

SBSeg 2008, Gramado, RS, Brazil

Information Leakage

- But C and D are internal cipher values !!
- Thus, we have discovered the input and output values of both f32 mappings of 2-round FOX64
- Consequently, we can attack each f32 mapping separately, and discover the round subkey of each f32 in turn.
- Conclusion: we can break 2-round FOX64

SBSeg 2008, Gramado, RS, Brazil

Information Leakage

- A similar reasoning applies to FOX128/k/2.
- Analogously, a similar approach holds to more than 2 rounds. But, rather than individual 32-bit values, we obtain an xor of 32-bit internal data.
- A similar kind of leakage was already observed in DES (due to D.W.Davies and S. Murphy), due to the Feistel structure.

SBSeg 2008, Gramado, RS, Brazil

Impossible-Differential Technique

- Developed by L.R.Knudsen and first applied on DEAL block cipher (1998)
- Biham also mentions its use (implicitly) in Enigma machine
- Impossible Diff. (ID) technique uses tools from differential cryptanalysis (DC)
- Unlike DC, though, ID uses differentials with probability zero (never hold)

SBSeg 2008, Gramado, RS, Brazil

Impossible-Differential Technique

- ID distinguishers consist of two (truncated) differentials, called Δ and ∇ , both of which hold with certainty.
- ∇ propagates differences in the encryption direction; Δ propagates differences in the decryption direction.
- The combination of Δ and ∇ makes up the ID distinguisher: $\nabla \not\rightarrow \Delta$ and $\nabla \not\rightarrow \Delta$.
- That is the reason for the name “impossible differential”.

SBSeg 2008, Gramado, RS, Brazil

Impossible-Differential Technique

- For FOX64, we have found a 3-round ID distinguisher.
- In particular, we have found that

$$(\Delta, \Delta, \Delta, \Delta) \not\rightarrow (\Delta, \Delta, 0, \Delta)$$
 where each Δ stands for a 32-bit nonzero xor difference.
- This fact by itself already allows one to distinguish 3-round FOX64 from a random permutation. But further, we can recover subkeys around this distinguisher.

SBSeg 2008, Gramado, RS, Brazil

Impossible-Differential Technique

- Attack complexities:
time: 2^{118} 5-round FOX64 computations
data: 2^{36} CP
memory: 2^{122} text blocks

time: 2^{198} 6-round FOX128 computations
data: 2^{52} CP
memory: 2^{249} text blocks
See attack details in the paper, please.

SBSeg 2008, Gramado, RS, Brazil

Conclusions and Open Problems

- This paper presented some cryptanalytic results of reduced-round FOX ciphers, such as information leakage and impossible differential attacks.
- The attacks do not threaten any official version of FOX (with the suggested parameters), but show new avenues for attacks.

SBSeg 2008, Gramado, RS, Brazil

Questions???

Send your doubts to

jorge_nakahara@yahoo.com.br

SBSeg 2008, Gramado, RS, Brazil